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A method for processing a database query according to at least one grouping column value, the method comprising:

partially pre-aggregating records in a database to provide a result that contains at least two records having like grouping column values; and

aggregating records derived from the partial pre-aggregation to provide a result that contains records having unique grouping column values.

2. The method as recited in claim 1, wherein the partially preaggregating further comprises:

maintaining a record store in memory, the record store having one record for each different grouping column value encountered in the operation;

receiving a new records

combining the new record with a record having the same grouping column value, if such a record exists; and

adding the new record to the record store in the memory if there is no record in the record store that has the same grouping column value as the new record.

3. The method as recited in claim 2, further comprising:

adding additional new records to the record store until the record store reaches a capacity such that it can accept no new records; and

outputting one or more records from the record store to a subsequent database operator.

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- The method as recited in claim 3, wherein after the one or more records have been output to the subsequent database operator, the adding and outputting are repeated until there are no new records to process.
- 5. The method as recited in claim 4, wherein any records remaining in the record store after there are no new records to process are output to the subsequent database operator.
- The method as recited in claim 3, wherein the subsequent database 6. operator is a join.
- 7. The method as recited in claim 1, further comprising estimating the costs and benefits of the partial pre-aggregation, and partially pre-aggregating the records only if the estimating indicates that the benefits are greater than the costs.
- 8. The method as recited in claim 1, wherein the partially preaggregating includes utilizing a hashing function.
- The method as recited in claim \(\), wherein the partial pre-aggregating 9. creates a record store in memory, and wherein the method further comprises utilizing the record store in memory for one or more other database operators.

10. One or more computer-readable media having computer-executable
instructions that, when executed by a computer, perform the method recited in
claim 1.
11. A computer programmed to perform the method recited in claim 1.
12. A relational database system, comprising:
memory for storing a record store, the memory having a portion available
for query processing;
a query processor coupled to the memory to process a query on the record *
store according to at least one grouping column value, the query processor being
configured to partially pre-aggregate the record store to provide a result that
contains at least two data records that have like grouping column values; and
the guery processor being further configured to aggregate data records

the query processor being further configured to aggregate data records resulting from the partial pre-aggregation to provide an aggregation result that contains data records, no two of the data records having a same grouping column value.

13. The relational database system as recited in claim 12, wherein the query processor being configured to partially pre-aggregate the record store comprises the query processor being configured to:

maintain a record store in the volatile memory, the record store having one record for each different grouping column value encountered in the partial preaggregation;

receive an input record from the non-volatile memory;

combine the input record with a record in the record store that has the same grouping column value, if there is such a record; and

adding the input record to the record store if there is no record in the record store that has the same grouping column value as the input record.

14. The relational database system as recited in claim 13, wherein the query processor is further configured to:

add additional input records to the record store in memory until the record store reaches a capacity such that it can accept no more input records; and output the records in the record store to a subsequent database operator.

15. The relational database system as recited in claim 14, wherein the query processor is configured to:

continually add input records to the record store; and

output one or more records from the record store to a subsequent database operator when the record store reaches a capacity such that it can accept no more new records, or whenever there are no new records to process.

- 16. The relational database system as recited in claim 12, wherein the query processor is further configured to perform a join on the records resulting from the partial pre-aggregation.
- 17. The relational database system as recited in claim 12, wherein the query processor is further configured to create a record store as a result of the partial pre-aggregation and utilize the record store in processing of another database operator.
- 18. The relational database system as recited in claim 12, further comprising:

a query optimizer configured to estimate the costs and benefits of the query processor performing a partial pre-aggregation; and

wherein the query processor performs the partial pre-aggregation only if the query optimizer indicates that the benefits of the partial pre-aggregation are greater than the costs of the partial pre-aggregation.

19. The relational database system as recited in claim 12, wherein the query processor is further configured to utilize hashing to perform the partial preaggregation.



'	20.	The	relational	database	system	as	recited	in	claim	12,	wherein	the
query	process	or is	further co	nfigured	to utiliz	e ha	ashing a	nd j	partitio	onin	g to perf	form
he pa	rtial pre	-agg	regation.									

21. A relational database computer program embodied on a computer-readable medium, comprising:

partial pre-aggregation code to partially pre-aggregate data records according to at least one grouping column value to provide a partial pre-aggregation result having two or more records having like grouping column values; and

aggregation code to aggregate data records in the partial pre-aggregation result to provide an aggregation result having records with unique grouping column values.

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The relational database computer program as recited in claim 21, wherein the partial pre-aggregation code is designed to:

maintain a record store in memory, the record store having one record for each different grouping column value encountered in the partial pre-aggregation;

receive à new record;

combine the new record with a record in the record store having the same grouping column value, if such a record exists;

add the new record to the record store no record in the record store that has the same grouping column value as the new record;

continuously add additional new records to the record store until the record store has reached a record store capacity;

output one or more records from the record store to a subsequent database operator when the record store has reached the record store capacity; and

output the records in the record store when there are no new records to process.

The relational database computer program as recited in claim 22, 23. further comprising database operator code that utilizes the record store for input.

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A relational database computer program stored on a computerreadable medium, the relational database computer program comprising computerexecutable instructions that, when executed on a computer, perform the following steps:

receiving\a stream of input records;

aggregating each input record in the stream as it is received to create a record store;

joining records in the record store with other data; and aggregating the records output from the join.

25. The relational database computer program as recited in claim 24, wherein:

the record store has a capacity that is less than the number of records in the stream of input records; and

the aggregating each input record is performed until the record store reaches capacity.

The relational database computer program as recited in claim 24, 26. further comprising computer-executable instructions that, when executed by a computer, perform the following steps:

determining, if it is optimal to aggregate the input records prior to performing the join; and

performing the aggregation prior to the join only if a determination is made that it is optimal to perform an aggregation prior to the join.